

A Mathematical Model for the Rainbow Vertex-Connection Number

Fidan Nuriyeva¹, Onur Ugurlu² and Hakan Kutucu³

¹ Institute of Cybernetics, Azerbaijan Academy of Sciences
Baku, Azerbaijan
nuriyevafidan@gmail.com

² Department of Mathematics, Ege University
Izmir, Turkey
onur_ugurlu@hotmail.com

³ Department of Computer Engineering, Karabuk University
Karabuk, Turkey
hakankutucu@karabuk.edu.tr

Abstract

The concept of rainbow connection was introduced by Chartrand et al. [2] in 2008. A vertex-colored graph is rainbow vertex-connected if any two vertices are connected by a path whose internal vertices have distinct colors. The rainbow vertex-connection of a connected graph G , denoted by $rvc(G)$, is the minimum number of colors needed to make G rainbow vertex-connected. In this paper, we introduce the first mathematical model of the problem.

Keywords: *Rainbow Vertex Connection, Graph Coloring, Integer Programming, Mathematical Modeling.*

1. Introduction

An edge coloring of a graph is a function from its edge set to the set of natural numbers. A path in an edge colored graph with no two edges sharing the same color is called a rainbow path. An edge colored graph is said to be rainbow connected if every pair of vertices is connected by at least one rainbow path. Such a coloring is called a rainbow coloring of the graph. The minimum number of colors required to rainbow color a connected graph G is called its rainbow connection number, denoted by $rc(G)$. For a basic introduction to the topic, see Chapter 11 in [3].

The concept of rainbow connectivity was recently introduced by Chartrand et al. in [2] as a measure of strengthening connectivity. The rainbow connection problem, apart from being an interesting combinatorial property, also finds an application in routing messages [4]. In addition to being a natural combinatorial measure, rainbow connectivity can be motivated by its interesting interpretation in the area of networking. Suppose that G represents a network (e.g., a cellular network). We wish to

route messages between any two vertices in a pipeline, and require that each link on the route between the vertices (namely, each edge on the path) is assigned a distinct channel (e.g. a distinct frequency). Clearly, we want to minimize the number of distinct channels that we use in our network. This number is precisely $rc(G)$ [11].

As one can see, the above rainbow connection number involves edge colorings of graphs. A natural idea is to generalize it to a concept that involves vertex-colorings. In [8], Krivelevich and Yuster are the first to introduce a new parameter corresponding to the rainbow connection number which is defined on a vertex-colored graph. A vertex-colored graph G is rainbow vertex-connected if its every two distinct vertices are connected by a path whose internal vertices have distinct colors. A vertex-coloring under which G is rainbow vertex-connected is called a rainbow vertex-coloring. The rainbow vertex-connection number of a connected graph G , denoted by $rcv(G)$ is the smallest number of colors that are needed in order to make G rainbow vertex-connected. The minimum rainbow vertex-coloring is defined similarly [13].

Rainbow connectivity from a computational point of view was first studied by Caro et al. [4] who conjectured that computing the rainbow connection number of a given graph is NP-Hard. This conjecture was confirmed by Chakraborty et al. [11], who proved that even deciding whether rainbow connection number of a graph equals 2 is NP-Complete. They also conjectured that for every integer $k \geq 2$ to decide whether $rc(G) \leq k$ is NP-Hard. Recently, Ananth and Nasre confirmed the conjecture in [10]. Li and Li [9] showed that for any fixed integer $k \geq 2$ to decide

whether $rc(G) \leq k$ is actually NP-Complete. For the rainbow vertex-connection number, we have a similar complexity result in [5].

The authors in [4], [11], [8] view rainbow connection number as a "quantifiable" way of strengthening the connectivity property of a graph. Hence, tighter upper bounds on the rainbow connection number for a graph with higher connectivity have been a subject of investigation. There are a number of algorithms that are proposed for the problem. Some of these are Basavaraju's approximation algorithm [1], Deng's polynomial algorithm [6], and Ugurlu and Nuriyeva's heuristic algorithm [12].

To the best of our knowledge, no mathematical model has been proposed for the rainbow vertex-connection problem. In this paper, we present a mathematical model of the rainbow vertex-connection problem.

2. Mathematical Modeling of the Rainbow Vertex-Connection Problem

The graphs we consider are finite, simple, undirected, and loopless. Let $G = (V, E)$ be a graph, where $V = \{v_1, v_2, \dots, v_n\}$ is the set of vertices, and $E = \{(v_i, v_j) | v_i, v_j \in V\}$ is the set of edges in G such that $|V| = n$ and $|E| = m$.

Let $A = a_{ij}$ be the adjacency matrix of graph G as follows:

$$a_{ij} = \begin{cases} 1, & \text{if there is an edge between } i \text{ and } j, \\ 0, & \text{otherwise} \end{cases}$$

We define the variables y_i and x_{ij}^{pq} as follows:

$$y_i \in \mathbb{N}$$

and

$$x_{ij}^{pq} = \begin{cases} 1, & \text{if the edge } (v_i, v_j) \text{ is on the } v_p \rightarrow v_q, \\ 0, & \text{otherwise,} \end{cases}$$

for all $(v_i, v_j) \in E$, $p = 1, \dots, n-1$, $q = p+1, \dots, n$. The variable y_i represents the color of the vertex $v_i \in V$.

Let v_p and v_q be the source vertex and the target vertex, respectively. The paths between all p, q pairs can be modeled as follows: Since there can be only one outgoing edge in a path from the source,

$$\sum_{i=1}^n x_{pi}^{pq} = 1, \\ \text{for all } (v_p, v_i) \in E, p = 1, \dots, n-1, q = p+1, \dots, n,$$

and there can be only one incoming edge in a path to the target,

$$\sum_{i=1}^n x_{iq}^{pq} = 1, \\ \text{for all } (v_i, v_q) \in E, p = 1, \dots, n-1, q = p+1, \dots, n.$$

Since the number of the incoming and outgoing edges are equal for each vertex in a path apart from the source and the target,

$$\sum_{i=1}^n x_{ik}^{pq} = \sum_{i=1}^n x_{ki}^{pq}, \\ \text{for all } (v_i, v_k) \in E, k = 1, \dots, n, p = 1, \dots, n-1, q = p+1, \dots, n: k \neq p, q.$$

Theorem 1 Since the internal vertices in a rainbow path connecting vertices v_p and v_q must be colored with different color (v_p and v_q may be colored or not colored with the same color), the following inequality is necessary for the vertex-rainbow problem:

$$\sum_{i=1}^{n-1} \sum_{j=1}^n (y_i + y_j) x_{ij}^{pq} \geq 2 \left(1 + 2 + \dots + \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} - 1 \right) \right) \\ \text{for all } (v_i, v_j) \in E$$

for all $p = 1, \dots, n-1$ and $q = p+1, \dots, n$.

Proof: It is clear that the number of edges in a rainbow path connecting vertices v_p and v_q is $\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq}$.

Therefore, in this rainbow path, the number of vertices

other than v_p and v_q is $\left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right) - 1$. Since source

and destination vertices in each rainbow path may not be colored and the other vertices in a rainbow path must be colored with different color, it is needed

$\left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right) - 1$ different colors to color these internal

vertices. Let y_i and y_j be the colors of the endpoints of an edge (v_i, v_j) . Then, the following constraint must be satisfied for a rainbow path connecting vertices v_p and v_q :

$$\sum_{i=1}^{n-1} \sum_{j=1}^n (y_i + y_j) x_{ij}^{pq} \geq 2 \left(1 + 2 + \dots + \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right) - 1 \right).$$

The constraint above can be re-written as follows:

$$\sum_{i=1}^{n-1} \sum_{j=1}^n (y_i + y_j) x_{ij}^{pq} \geq \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right) - 1 \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right).$$

Consequently, the rainbow vertex-connection problem can be formulated as the following integer nonlinear program:

$$rvc(G) = \min \sum_{i=1}^n y_i \quad (1)$$

subject to

$$\sum_{i=1}^n x_{pi}^{pq} = 1, \quad (2)$$

for all $(v_p, v_i) \in E$, $p = 1, \dots, n-1$, $q = p+1, \dots, n$.

$$\sum_{i=1}^n x_{iq}^{pq} = 1, \quad (3)$$

for all $(v_i, v_q) \in E$, $p = 1, \dots, n-1$, $q = p+1, \dots, n$.

$$\sum_{i=1}^n x_{ik}^{pq} = \sum_{i=1}^n x_{ki}^{pq} \quad (v_k, v_i) \in E \quad (4)$$

for all $k = 1, \dots, n$, $p = 1, \dots, n-1$, $q = p+1, \dots, n$; $k \neq p, q$.

$$\sum_{i=1}^{n-1} \sum_{j=1}^n (y_i + y_j) x_{ij}^{pq} \geq \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right) - 1 \left(\sum_{i=1}^{n-1} \sum_{j=1}^n x_{ij}^{pq} \right).$$

for all $p = 1, \dots, n-1$ and $q = p+1, \dots, n$. (5)

$$x_{ij}^{pq} = 0 \vee 1, \text{ for all } (v_i, v_j) \in E, p = 1, \dots, n-1, q = p+1, \dots, n. \quad (6)$$

$$y_i = k \text{ for all } (v_i, v_j) \in E, \text{ where } k \in \mathbb{N}. \quad (7)$$

The proposed model satisfies the rainbow coloring of the graphs and $\max\{y_i \mid v_i \in V\}$ gives the minimum rainbow vertex-connection number. The number of variables and constraints in the model are $O(mn^2)$ and $O(n^3)$, respectively.

3. Conclusions

The rainbow vertex-connection number is a new concept for measuring the connectivity of a graph. In this paper, we present a mathematical model of the rainbow vertex-connection problem. In future works, we aim to investigate the computational results on graphs with various densities.

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Fidan Nuriyeva is a Senior Researcher received her B.Sc. degree in Mathematics and Computer Science in 2008 and M.Sc. degree in Computer Science from Ege University in 2010, received her Ph.D. degree in Computer Science from Ege University in 2013. She is currently at the Institute of Cybernetics, Azerbaijan Academy of Sciences in Baku. Her research interests include combinatorics and optimization algorithms.

Onur Ugurlu received his B.Sc. degree in Mathematics and Computer Science in 2011 and M.Sc. degree in Computer Science from Ege University in 2013. He is currently a Ph.D. candidate in Computer Science at Ege University, in Turkey. His research interests include optimization, graph theory and design and analysis of algorithms.

Hakan Kutucu is an assistant professor in the department of computer engineering at Karabuk University in Turkey. He has two M.Sc. degrees, in mathematics and in computer engineering, and a Ph.D. in mathematics & computer science. His primary research area is network design problems, mathematical modeling and linear programming.